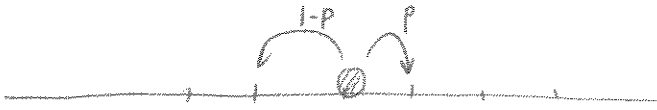


examples

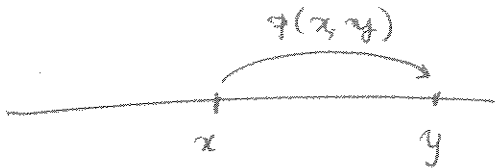
1) Random walk on \mathbb{Z} (Simple)



$$P(i, i+1) = p \quad i \in \mathbb{Z}$$

$$P(i, i-1) = 1-p$$

2) Random walk on \mathbb{Z}



Define $p(x,y) = P\{X_1=y | X_0=x\}$

$$\sum_{y \in S} p(x,y) = 1$$

$S = \mathbb{Z}$, \Rightarrow infinite sum.

n -step transitions : $P_n(x,y) = P\{X_n=y | X_0=x\}$

No matrix ?!

Chapman - Kolmogorov equations

$$\begin{aligned} \text{If } 0 < m, n < \infty \quad P_{m+n}(x,y) &= P\{X_{m+n}=y | X_0=x\} \\ &= \sum_{z \in \mathbb{Z}} P\{X_{m+n}=y, X_m=z | X_0=x\} \\ &= \sum_{z \in \mathbb{Z}} P\{X_{m+n}=y | X_m=z\} P\{X_m=z | X_0=x\} \\ &= \sum_{z \in \mathbb{Z}} P_m(x,z) \cdot P_n(z,y). \end{aligned}$$

\rightarrow $P_{m+n}(x,y) = \sum_{z \in \mathbb{Z}} P_m(x,z) P_n(z,y)$ This should make sense!

3) Random walk with partially reflecting boundary.

$0 < p < 1$, $S = \{0, 1, 2, \dots\}$



$P(x, x+1) = p$

$P(x, x-1) = 1-p$ $x \neq 0$

$P(0, 0) = 1-p$

Queuing model. (Discrete time)

Let $X_n = \#$ of customers waiting in line for some service.

During each time interval, with prob. p someone new arrives

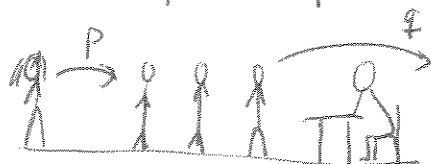
With probability q someone gets serviced and leaves the queue.

- This is a Markov Chain. , $S = \{0, 1, 2, \dots\}$

- Transition probabilities:

$P(x, x-1) = q(1-p)$ (why?) $P(x, x) = qP + (1-q)(1-p)$ $x \neq 0$
 $P(x, x+1) = p(1-q)$ $x \neq 0$

$P(0, 0) = 1-p$, $P(0, 1) = p$



Questions : Can the line become infinite?

Is there a stationary distribution?

Section 2.2. Recurrence and Transience.

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n is an irreducible MC, countably infinite state space S and transition probabilities $p(x, y)$

Definitions:

The chain X_n is recurrent if for each state $x \in S$,

$$\mathbb{P}\{X_n = x \text{ for infinitely many } n\} = 1$$

The chain X_n is transient if for each state $x \in S$,

$$\mathbb{P}\{X_n = x \text{ for infinitely many } n\} = 0.$$

Theorem.

Lec. 16

An irreducible MC is transient if and only if the expected number of returns to a state x is finite, i.e. iff

$$\sum_{n=0}^{\infty} P_n(x, x) < \infty, \text{ for some } x.$$

Proof WLOG assume $X_0 = x$.

Let $R_x = \sum_{n=0}^{\infty} \mathbb{1}\{X_n = x\}$ = # of total visits to state x .

If $R_x \neq \infty$ then the chain cannot be recurrent, so it's transient.

But if $R_x < \infty$ then we can actually compute its expectation,

$$\begin{aligned} \mathbb{E}(R_x) &= \mathbb{E}\left(\sum_{n=0}^{\infty} \mathbb{1}\{X_n = x\}\right) = \sum_{n=0}^{\infty} \mathbb{E}(\mathbb{1}\{X_n = x\}) \\ &= \sum_{n=0}^{\infty} \mathbb{P}\{X_n = x\} \stackrel{\text{why?}}{=} \sum_{n=0}^{\infty} P_n(x, x). \end{aligned}$$

Now ~~can~~ define

$$T_x^1 = \min\{n > 0 : X_n = x\}, \text{ the time of first return to } x.$$

We say $T_x = \infty$ if the chain does not return to x .

Assume by way of contradiction that $\mathbb{P}\{T_x^1 < \infty\} = 1$ (*)

(so we always return).

Recurisively, define the time of k -th return

$$T_x^k = \min \{n > T_x^{k-1} : X_n = x\}, k \geq 2$$

As discussed before how $T_x^k - T_x^{k-1} \stackrel{D}{=} T_x^1$ and that

$$T_x^k - T_x^{k-1} \perp T_x^j - T_x^{j-1} \quad \forall j \neq k.$$

Thus, $\mathbb{P}\{T_x^k - T_x^{k-1} < \infty\} = 1 \quad \forall k.$

But, $\mathbb{P}\{\text{we visit state } x \text{ infinitely many times}\} =$ ~~1~~

$$= \lim_{k \rightarrow \infty} \mathbb{P}\{\text{we visit state } x \text{ } k \text{ times}\}$$

$$= \lim_{k \rightarrow \infty} \mathbb{P}\{T_x^k < \infty\} = \lim_{k \rightarrow \infty} \mathbb{P}\{T_x^k - T_x^{k-1} + T_x^{k-1} - T_x^{k-2} + \dots + T_x^1 < \infty\}$$

$$\stackrel{\text{why?}}{=} \lim_{k \rightarrow \infty} \mathbb{P}\{T_x^k - T_x^{k-1} < \infty, T_x^{k-1} - T_x^{k-2} < \infty, \dots, T_x^1 < \infty\}$$

$$\stackrel{\text{why?}}{=} \lim_{k \rightarrow \infty} \mathbb{P}\{T_x^k - T_x^{k-1} < \infty\}^k = 1.$$

contradiction.
↑ this shows that (*) ⇔ recurrence

Now assume that $\mathbb{P}\{T_x^1 < \infty\} = q < 1$

Let's compute $\mathbb{P}\{R_x = m\}$.

$$\cdot \mathbb{P}\{R_x = 1\} = \mathbb{P}\{\text{chain never returns}\} = 1 - q \quad (\text{why?})$$

$$\cdot \mathbb{P}\{R_x = 2\} = \mathbb{P}\{\text{chain returns once and never again}\} = \mathbb{P}\{T_x^1 < \infty, T_x^2 - T_x^1 = \infty\} = q(1-q) \quad (\text{why?})$$

Claim $P\{R_x = m\} = q^{m-1}(1-q)$ (Geometric! Makes sense?)

$$\Rightarrow E(R_x) = \sum_{m=1}^{\infty} m P\{R_x = m\} = \sum_{m=1}^{\infty} m (1-q) q^{m-1}$$

$$= \frac{1}{1-q} < \infty$$

$$\Rightarrow \boxed{\frac{1}{1-q} = E(R_x) = \sum_{n=0}^{\infty} P_n(x, x) < \infty.}$$

By-product formula:

$$\boxed{\sum_{n=0}^{\infty} P_n(x, x) = \frac{1}{1 - P\{T_x^{\pm} < \infty\}} = E(R_x)}$$

Examples

Simple symmetric RW in \mathbb{Z}^d : Determine whether are recurrent or transient

Case 1. $d=1$

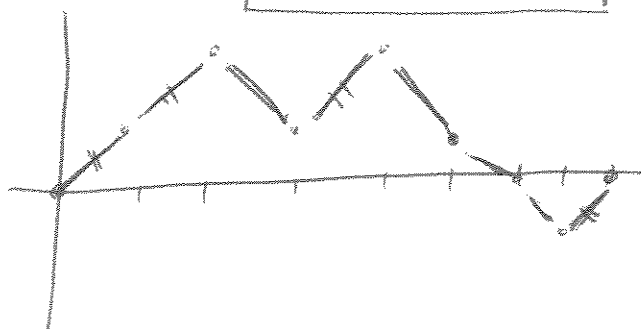
$p(x, x+1) = p(x, x-1) = \frac{1}{2}$. Note: This is a periodic chain!

Because it is irreducible, we need to focus on just one state, say 0 and assume that $X_0 = 0$.

We want to estimate the sum $\sum_{n=0}^{\infty} P_n(0, 0)$.

First note that $P_n(0, 0) = 0$ if n is odd! So we need

to evaluate $\sum_{n=0}^{\infty} P_{2n}(0, 0)$.



For the chain to return to 0, the walk took an equal number of left and right steps, and that number must be n .

Thus $P_{2n}(0,0) = \binom{2n}{n} \cdot \left(\frac{1}{2}\right)^{2n} = \frac{(2n)!}{n!n!} \left(\frac{1}{2}\right)^{2n}$.

We need to estimate $n!$ to manipulate the sum. Use Taylor's

formula (proof later) $n! \sim \sqrt{2\pi n} n^n e^{-n}$

So $P_{2n}(0,0) = \frac{(2n)!}{n!n!} \frac{1}{(2)^{2n}} \approx \frac{\sqrt{4\pi n} 2^{2n} n^{2n} e^{-2n}}{2\pi n n^n n^n e^{-n} e^{-n}} \cdot \frac{1}{2^{2n}}$

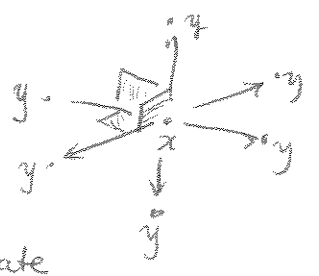
$= \frac{1}{\sqrt{\pi n}}$
 $\Rightarrow \sum_{n=1}^{\infty} P_{2n}(0,0) \approx \sum_{n=1}^{\infty} \frac{1}{\sqrt{\pi n}} = +\infty$ (why?).

So 1-dimensional r.w. simple symm. is recurrent.

Lec. 17

use 2. d=3

Transition probabilities: $p(x,y) = \begin{cases} \frac{1}{6} & \text{if } x \sim y \\ 0 & \text{o.w.} \end{cases}$



Again, for n odd $P_n(0,0) = 0$, so we need to estimate

$\sum_{n=0}^{\infty} P_{2n}(0,0)$.

As before, the walk must have taken equal # steps in each of the directions, say i, j, k , with $i+j+k=n$.

$\Rightarrow P_{2n}(0,0) = \sum_{\substack{i,j,k \geq 0 \\ i+j+k=n}} \frac{(2n)!}{(i!j!k!)^2} \left(\frac{1}{6}\right)^{2n} = \sum_{\substack{i,j,k \geq 0 \\ i+j+k=n}} \frac{(2n)!}{n!n!} \cdot \frac{n!n!}{(i!j!k!)^2} \cdot \left(\frac{1}{6}\right)^{2n}$
 $= \binom{2n}{n} \left(\frac{1}{2}\right)^{2n} \cdot \sum_{\substack{i,j,k \geq 0 \\ i+j+k=n}} \binom{n}{i,j,k}^2 \cdot \left(\frac{1}{3^2}\right)^{2n}$ (*)

if $n = 3m$ then $\binom{n}{i, j, k} \stackrel{(\dagger)}{\leq} \binom{3m}{m, m, m}$ (55)

$$\begin{aligned} \Rightarrow (*) &\leq \binom{2n}{n} \left(\frac{1}{2}\right)^{2n} \sum_{\substack{i, j, k \geq 0 \\ i+j+k=n}} \binom{n}{i, j, k} \binom{3m}{m, m, m} \frac{1}{3^n} \cdot \frac{1}{3^n} \\ &= \binom{2n}{n} \left(\frac{1}{2}\right)^{2n} \left(\frac{1}{3}\right)^n \binom{3m}{m, m, m} \underbrace{\sum \binom{n}{i, j, k} \frac{1}{3^n}}_{=1 \text{ (why?)}} \\ &\sim C \left(\frac{6}{n}\right)^{3/2}. \end{aligned}$$

(†) $\frac{n!}{(i-1)!(j+1)!k!} = \frac{n!}{i!j!k!} \cdot \frac{i}{j+1}$ if $i < j$
 so highest value is when i, j, k are equal. Hence the treatment of the case $n=3m$ separately.

f $n = 3m-1, 3m-2 \rightarrow P_{6m}^{(0,0)} \geq \frac{1}{6^2} P_{6m-2}^{(0,0)}$
 $P_{6m}^{(0,0)} \geq \frac{1}{6^4} P_{6m-4}^{(0,0)}$

$$\Rightarrow \sum P_n^{(0,0)} < C \sum \left(\frac{6}{n}\right)^{3/2} < \infty \text{ (why?)}$$

So 3-dimensional r.w. is transient!

Theorem.

Simple r.w. in \mathbb{Z}^d is recurrent if $d=1$ or 2 and transient

if $d \geq 3$. ●

Exercise 2.18. (Stirling's formula)

Let $\{X_i\}_{1 \leq i < \infty}$ an i.i.d. sequence of Poisson (1) r.v.

$$\left(\Rightarrow \mathbb{P}\{X_i = k\} = \frac{e^{-1}}{k!} \right)$$

Then $\sum_{i=1}^n X_i = Y_n \sim \text{Poisson}(n)$ r.v.

Let $p(n, k) = P\{Y_n = k\} = e^{-n} \frac{n^k}{k!}$ and let $a > 0$.

By CLT

$$P\left\{0 \leq \frac{Y_n - n}{\sqrt{n}} \leq a\right\} = P\left\{0 \leq \frac{\sum X_i - n}{\sqrt{n}} \leq a\right\}$$

$$\xrightarrow{n \rightarrow \infty} \int_0^a \frac{1}{\sqrt{2\pi}} e^{-x^2/2} dx.$$

But,

$$P\left\{0 \leq \frac{Y_n - n}{\sqrt{n}} \leq a\right\} = \sum_{k=n}^{n+a\sqrt{n}} P\{Y_n = k\} = \sum_{k=n}^{n+a\sqrt{n}} p(n, k).$$

Hence,

$$\lim_{n \rightarrow \infty} \sum_{k=n}^{n+a\sqrt{n}} p(n, k) = \int_0^a \frac{1}{\sqrt{2\pi}} e^{-x^2/2} dx. \quad (1)$$

Now,

$$p(n, k) = e^{-n} \frac{n^k}{k!} \quad \text{and assume} \quad n \leq k < n+a\sqrt{n}.$$

Then $\frac{n!}{k!} \leq n^{n-k}$ since we can do induction,

$$\Rightarrow \frac{n^k}{k!} \leq \frac{n^n}{n!} \Rightarrow e^{-n} \frac{n^k}{k!} \leq \frac{n^n}{n!} e^{-n} = p(n, n). \quad (2)$$

Also for n large

$$e^{-a^2} \leq \left(1 - \frac{a}{\sqrt{n}+a}\right)^{a\sqrt{n}} \leq \left(\frac{n}{n+a\sqrt{n}}\right)^{k-n} \leq \frac{n^{k-n}}{k^{k-n}} \leq \frac{n^{k-n} n!}{k!}$$

$$\Rightarrow e^{-a^2} p(n, n) \leq p(n, k) \quad (3)$$

From (1), (2), (3)

$$\Rightarrow p(n, n) \sim \frac{1}{\sqrt{2\pi n}}$$

(Need to let $a \rightarrow 0$)

Define $\alpha(z) = \mathbb{P}\{X_n = z \text{ for some } n \geq 0 \mid X_0 = x\}$.

$\alpha(z) = 1$

If chain is recurrent, then $\alpha(z) = 1 \quad \forall z$!

If chain is transient, then there must be states with $\alpha(x) < 1$.
(makes sense?)

In fact, there should be states x further and further away from z with $\alpha(x)$ as small as we want.

$$\begin{aligned} \text{Also; } z \neq x, \alpha(x) &= \mathbb{P}\{X_n = z \text{ for some } n \geq 0 \mid X_0 = x\} \\ &= \mathbb{P}\{X_n = z \text{ for some } n \geq 1 \mid X_0 = x\} \\ &= \sum_{y \in S} \mathbb{P}\{X_1 = y \mid X_0 = x\} \mathbb{P}\{X_n = z \text{ for some } n \geq 1 \mid X_1 = y\} \\ &= \sum_{y \in S} p(x, y) \alpha(y). \end{aligned}$$

For a transient chain, $\alpha(x)$ satisfies

- (1) $0 \leq \alpha(x) \leq 1$
- (2) $\alpha(z) = 1, \quad \inf\{\alpha(x) : x \in S\} = 0$
- (3) $\alpha(x) = \sum_{y \in S} p(x, y) \alpha(y), \quad x \neq z.$

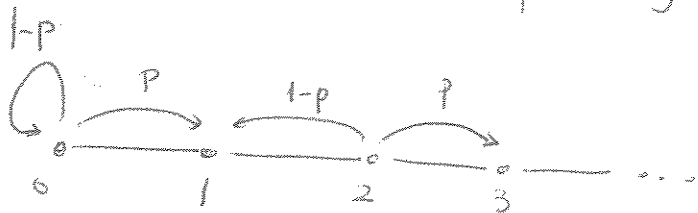
Criterion.

X_n recurrent \iff no solution $\alpha(x)$ satisfying (1), (2), (3) exists.

X_n transient $\iff \exists!$ unique solution to (1), (2), (3).

Example.

Random walk with partially reflecting boundary.



$$p(x, x-1) = 1-p \quad x \geq 1$$

$$p(x, x+1) = p \quad x \geq 0$$

$$p(0,0) = 1-p$$

Intuition?

Let $z=0$ and see if we can solve (1), (2), (3).

$$y \quad (3) \quad \alpha(x) = (1-p)\alpha(x-1) + p\alpha(x+1), \quad x > 0, \quad \alpha(0) = 1.$$

$$\implies \alpha(x) = c_1 + c_2 \left(\frac{1-p}{p}\right)^x \quad p \neq \frac{1}{2}$$

$$\alpha(x) = c_1 + c_2 \cdot x \quad p = \frac{1}{2}$$

We have that $\alpha(0) = 1 \implies$

$$\alpha(x) = \begin{cases} (1-c_2) + c_2 \left(\frac{1-p}{p}\right)^x & p \neq \frac{1}{2} \\ 1 + c_2 x & p = \frac{1}{2} \end{cases}$$

If $c_2 = 0$ we cannot satisfy (2) [since $\alpha(x) = 1 \forall x$]

If $c_2 \neq 0$, $p = \frac{1}{2}$, $\alpha(x) \notin (0,1) \forall x$, } \rightarrow (1) is violated.
 If $c_2 \neq 0$, $p < \frac{1}{2}$, $\alpha(x) \notin (0,1) \forall x$, }

So, no matter what c_2 is, (1), (2), (3) cannot be simultaneously satisfied, so the chain is recurrent for $p \leq \frac{1}{2}$

or $p > \frac{1}{2}$, $\inf_{x \in S} \{ \alpha(x) : x \in S \} = 0 \iff \inf_{x \in S} \{ (1-c_2) + c_2 \left(\frac{1-p}{p} \right)^x \} = 0$

$\iff \inf_{x \in S} \left\{ c_2 \left(\frac{1-p}{p} \right)^x \right\} = c_2 + 1 = 0$

$\iff -c_2 + 1 = 0$ (why?) , $1 = c_2 \implies \alpha(x) = \left(\frac{1-p}{p} \right)^x \in (0, 1)$

Also, $\alpha(0) = 0$ as required.

Since a solution to (1), (2), (3) exist, the chain is ~~recurrent!~~ transient!

Section 2.3 Positive & null Recurrence.

Question 1: X_n is irreducible, aperiodic, MC, on an infinite state space S

Does this imply the existence of a limiting (invariant) probability?

A1: No! X_n might be transient!

Q2: What if X_n is recurrent?

A2: Not always! \downarrow

A limiting probability distribution π satisfies

$\forall x, y \in S \quad \lim_{n \rightarrow \infty} P_n(y, x) = \pi(x)$

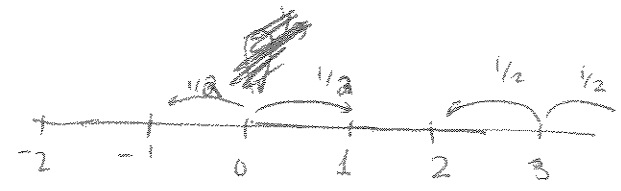
If X_n is transient then $\lim_{n \rightarrow \infty} P_n(y, x) = 0!$

\Rightarrow No limiting $\pi(x)$ exists ($\sum_{x \in S} \pi(x) \stackrel{!}{=} 1$ should.)

However, there's no guarantee if a chain is recurrent that

$$\lim_{n \rightarrow \infty} P_n(y, x) \neq 0.$$

For example, consider



This is irreducible, periodic, aperiodic.

~~$P_{2n+1}(0,0) > \frac{1}{3^{2n+1}} > 0$ (why?)~~

~~$P_{2n}(0,0) > \binom{2n}{n} \cdot \frac{1}{2^{2n-2}} \cdot \frac{1}{3^2}$ (why?)~~

Then show that $\sum_{n=0}^{\infty} P_n(0,0) = +\infty$, the chain is recurrent.

But $\lim_{n \rightarrow \infty} P_n(0,0) \rightarrow 0$ (~~show that~~ ~~periodic~~)

If you want to stay at 0 for k times, then $n-k$ has to be even (why?)

Then $P_n(0,0) \leq \sum_{k=0}^n \frac{1}{3^k} \binom{n-k}{\frac{n-k}{2}} \binom{n}{k} \frac{1}{2^{n-k}} \leq \binom{n}{\frac{n}{2}} \frac{1}{2^n} \sum_{k=0}^n \binom{2}{3}^k \binom{n}{k}$

Even though it's 2-periodic, it is recurrent but $P_{2n}(0,0) \rightarrow 0$.

We need to distinguish:

If X_n is an irreducible MC, aperiodic ~~with~~, recurrent with $P_n(x,y) \rightarrow 0 \Rightarrow$ Null recurrent.

O.W. $P_n(x,y) \rightarrow \pi(x) > 0 \Rightarrow$ Positive recurrent.

² Positive Recurrent Chains ^{Lec. 19} behave similarly to finite MC. (6)

Facts: If X_n irreducible, positive recurrent, aperiodic then

$$1) \lim_{n \rightarrow \infty} P_n(y, x) = \pi(x) > 0$$

$$2) \sum_{y \in S} \pi(y) p(y, x) = \pi(x) \quad (4) \text{ is an invariant distribution.}$$

3) The return time $T = \min\{n > 0 \mid X_n = x\}$ has

$$\mathbb{E}[T \mid X_0 = x] = 1/\pi(x)$$

If X_n is null recurrent, $\mathbb{E}(T_x) = +\infty$ but $T < \infty$ a.s.

The invariant distribution is unique.

If the chain is not positive recurrent, then no such π (as in (4)) exist.

Strategy: To check positive recurrence, try and find a distribution satisfying (4). If you show that no such π exist, you can't have positive recurrence.

Example: Random walk with partially reflecting boundary (again)

(4) gives

$$(1-p)\pi(x+1) + p\pi(x-1) = \pi(x) \quad x > 0.$$

$$\pi(1)(1-p) + \pi(0)(1-p) = \pi(0) \implies \pi(0) = \pi(1) \left[\frac{1-p}{p} \right].$$

The solution to the recursion is

$$\pi(x) = c_1 + c_2 \left(\frac{p}{1-p} \right)^x, \quad p \neq \frac{1}{2} \quad (a)$$

$$\pi(x) = c_1 + c_2 x, \quad p = \frac{1}{2} \quad (b)$$

$$\pi(0) = \frac{1-p}{p} \pi(1) \quad (*)$$

Plug (*) into (a) & (b), $\lim_{x \rightarrow \infty} \pi(x) = 0$ (why?)

$$\Rightarrow \pi(x) = c_2 \left(\frac{p}{1-p} \right)^x, \quad p \neq \frac{1}{2}$$

$$\pi(x) = c_1, \quad p = \frac{1}{2}$$

Need:

$$\sum_{x \in \mathbb{N}_0} \pi(x) = 1 \Rightarrow \text{For } p = \frac{1}{2} \text{ we have null recurrence}$$

$$\text{for } p < \frac{1}{2}, \quad 1 = c_2 \sum_x \left(\frac{p}{1-p} \right)^x = c_2 \cdot \frac{1}{1 - \frac{p}{1-p}} = c_2 \frac{1-p}{1-2p}$$

$$\Rightarrow c_2 = \frac{1-2p}{1-p}$$

$$\Rightarrow \boxed{\pi(x) = \left(\frac{1-2p}{1-p} \right) \left(\frac{p}{1-p} \right)^x} \quad (*)$$

Lec. 20

Discrete time, Countable state space MC:

For the purposes of the diagram, assume that we have an irreducible MC on $\mathbb{N} \cup \{0\}$. Of course, the diagram is true in more generality.

Irreducible MC (so we only need to consider one state)

Ways to show
Transient

1. $\sum_{n=0}^{\infty} P_n(i, i) < \infty$
for some state i

2. $\exists i$ s.t. $\mathbb{P}\{\text{ever returning to } i \mid X_0 = i\} < 1$

\Rightarrow
3. $\exists i$ s.t. $\mathbb{P}\{\text{Never returning to } i \mid X_0 = i\} > 0$

4. \exists a function $\alpha(x)$ s.t.

$$0 \leq \alpha(j) \leq 1$$

$$\alpha(i) = 1, \inf_{j \in S} \alpha(j) = 0$$

$$\alpha(j) = \sum_{k \in S} p(j, k) \alpha(k) \quad i \neq j \quad (*)$$

$$\alpha(j) = \mathbb{P}\{X_n = i \text{ for some } n \mid X_0 = j\}$$

Ways to show
Recurrence

1. $\sum_{n=0}^{\infty} P_n(i, i) = +\infty$
for some state i

2. $\exists i$ s.t. $\mathbb{P}\{\text{ever returning to } i \mid X_0 = i\} = 1$

\Leftrightarrow
3. $\exists i$ s.t. $\mathbb{P}\{\text{Never returning to } i \mid X_0 = i\} = 0$

4. No such function $\alpha(x)$ exists

Positive Recurrent

Null recurrent

a) Find a convenient state i to compute the expected time of first return

$$E[T_{i,i}^1] = \sum_{n=0}^{\infty} n \mathbb{P}\{T_{i,i}^1 = n\} < \infty$$

b)

Next time

a) Find a conv. state i to compute expected time of first return

$$E[T_{i,i}^1] = \infty$$

b)

Next time

Positive Recurrence

b) Find a function $\pi(x)$ s.t.

$$\left\{ \begin{array}{l} 0 \leq \pi(x) \leq 1 \\ \sum_{x \in S} \pi(x) = 1 \\ \pi(x) = \sum_{y \in S} \pi(y) p(y, x) \quad \forall x \in S \end{array} \right\}$$

π is the invariant distribution, and

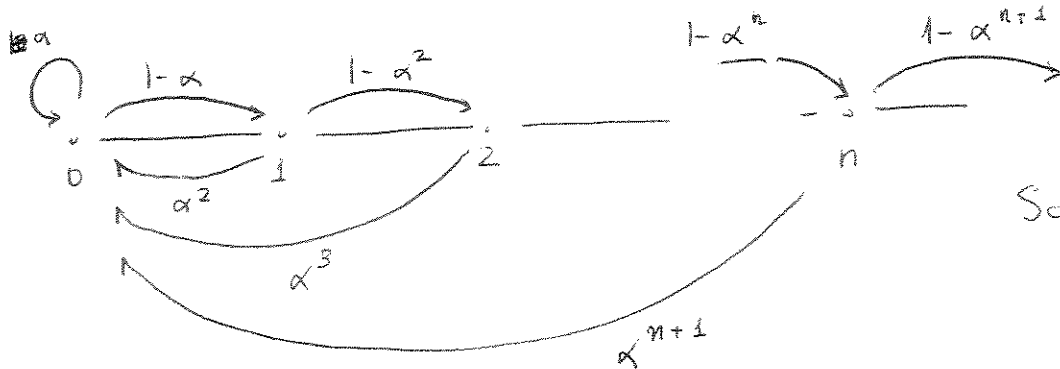
$$\pi(x) = \frac{1}{\mathbb{E}(T_{xx}^1)} > 0$$

Null Recurrence

b) Such function π cannot exist.

Examples

Let $0 < \alpha < 1$ and consider the MC (Give students ~7 minutes to work alone) then do on board



So $p(n, n+1) = 1 - \alpha^{n+1}$
 $p(n, 0) = \alpha^{n+1}$

1. Show / explain / check this is a periodic.
2. Check whether is recurrent or transient, using (3).

• Compute $\mathbb{P}\{\text{we never return to } 0 \mid X_0 = 0\} =$

$$= (1 - \alpha) \cdot (1 - \alpha^2) \cdot (1 - \alpha^3) \cdots (1 - \alpha^n) \cdots = \prod_{i=1}^{\infty} (1 - \alpha^i)$$

We say the infinite product converges iff $\lim_{n \rightarrow \infty} \prod_{i=1}^n (1 - \alpha^i)$ exists

and it is not equal to 0

A useful criterion for convergence of infinite products of this form:

$$\prod_{i=1}^{\infty} (1 \pm u_i) \text{ converges iff } \sum_{i=1}^{\infty} u_i < \infty$$

(Try and prove it by yourselves)

In this case, $u_i = \alpha^i$

$$\Rightarrow \text{Check } \sum_{i=1}^{\infty} \alpha^i = \frac{\alpha}{1-\alpha} < +\infty \text{ Hence the product}$$

converges and is positive. Thus,

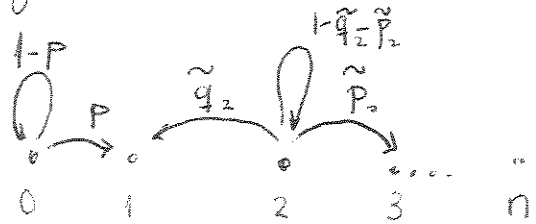
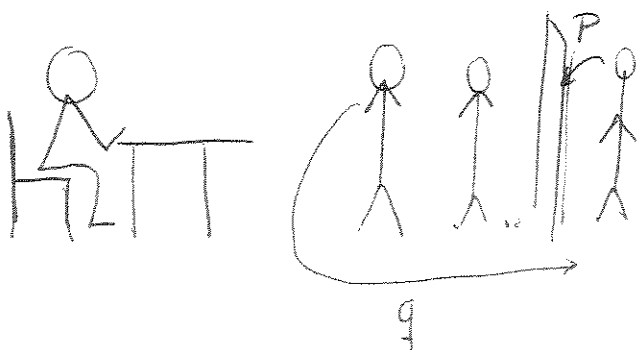
$\mathbb{P}\{\text{never return to } 0 \mid X_0 = 0\} > 0$, so the chain is transient. \blacksquare

Example 2: Single server discrete time queue.

Let X_n be the number of customers waiting to get serviced in a single queue. Let $0 < p, q < 1$

The probability of a customer arrival in a time window is p

The probability of a customer departure in a time window is q as long as there is a customer to get serviced!



$$\tilde{p}_2 = p(1-q) \text{ why? } = p(x, x+1) \quad x > 0 \quad p(0,0) = 1-p = \tilde{p}_0$$

$$\tilde{q}_2 = q(1-p) \text{ why? } = p(x, x-1) \quad x \geq 1 \quad p(0,1) = p = \tilde{q}_0$$

Assumption k customers can be served in the same time frame as k served

Recurrence / Transience for single server queue

We use (4).

$$\text{Let } \alpha(n) = \mathbb{P}\{\text{we reach } 0 \mid X_0 = n\} = \alpha_n$$

Then (either by one-step analysis) or by (*), we have

$$\alpha_n = \tilde{p}_n \alpha_{n+1} + \tilde{q}_n \alpha_{n-1} + (1 - \tilde{p}_n - \tilde{q}_n) \alpha_n, \quad n > 0 \iff$$

$$(\tilde{p}_n + \tilde{q}_n) \alpha_n = \tilde{p}_n \alpha_{n+1} + \tilde{q}_n \alpha_{n-1} \quad n \geq 1 \quad \begin{array}{l} \text{ALGEBRA} \\ \iff \end{array}$$

$$\alpha_n - \alpha_{n+1} = \frac{\tilde{q}_n}{\tilde{p}_n} [\alpha_{n-1} - \alpha_n] \quad n \geq 1.$$

Iterating yields

$$\alpha_n - \alpha_{n+1} = \frac{\tilde{q}_1 \tilde{q}_2 \dots \tilde{q}_n}{\tilde{p}_1 \tilde{p}_2 \dots \tilde{p}_n} [\alpha_0 - \alpha_1] \stackrel{!}{=} \frac{\tilde{q}_1 \dots \tilde{q}_n}{\tilde{p}_1 \dots \tilde{p}_n} [1 - \alpha_1].$$

why?

Then,

$$\begin{aligned} \alpha_{n+1} &= [\alpha_{n+1} - \alpha_0] + \alpha_0 = \sum_{k=0}^n (\alpha_{k+1} - \alpha_k) + 1 \\ &= 1 - [1 - \alpha_1] \sum_{k=0}^n \frac{\tilde{q}_1 \dots \tilde{q}_k}{\tilde{p}_1 \dots \tilde{p}_k} \\ &= 1 - C \sum_{k=0}^n \left(\frac{\tilde{q}}{\tilde{p}} \right)^k \quad (**) \end{aligned}$$

If the sum (**) diverges (as $n \rightarrow \infty$) then the first two equations in (4) cannot be satisfied, so the chain is ~~transient~~ recurrent. The sum converges only if $\tilde{p} > \tilde{q} \iff$

⇔

$$p(1-q) > q(1-p) \iff p > q$$

Lemma 11. If $p > q$ the chain is transient and the server should either be fired or get a medal

If $p \leq q$ the chain is recurrent

Question: When do you think we have null recurrence?

Invariant distributions.

We use ~~(a)~~ the set of equations in (b). Then

$$\pi(x) = \pi(x-1) \tilde{p}_{x-1} + \pi(x) (1 - \tilde{p}_x - \tilde{q}_x) + \pi(x+1) \tilde{q}_{x+1} \quad x \geq 1 \quad (+)$$

$$\pi(0) = \pi(0) (1 - \tilde{p}_0) + \pi(1) \tilde{q}_1 \quad (+')$$

$$\iff \tilde{q}_1 \pi(1) - \tilde{p}_0 \pi(0) = 0$$

From (+) algebra,

$$\tilde{q}_{x+1} \pi(x+1) - \tilde{p}_x \pi(x) = \tilde{q}_x \pi(x) - \tilde{p}_{x-1} \pi(x-1)$$

$$\begin{aligned} &= \dots \\ &= \tilde{q}_1 \pi(1) - \tilde{p}_0 \pi(0) \\ &= 0 \end{aligned}$$

$$\iff \boxed{\pi(x+1) = \frac{\tilde{p}_x}{\tilde{q}_{x+1}} \pi(x)} \quad \forall x \geq 0$$

$$\begin{aligned} &= \frac{\tilde{p}_0 \tilde{p}_1 \dots \tilde{p}_{x-1}}{\tilde{q}_1 \tilde{q}_2 \dots \tilde{q}_x} \pi(0) = \frac{\tilde{p}_0}{\tilde{p}_0} \left(\frac{\tilde{p}_x}{\tilde{q}_x} \right)^{x-1} \pi(0) \\ &= \left(\frac{\tilde{p}_x}{\tilde{q}_x} \right)^x \pi(0) \end{aligned}$$

Then, we want.

$$\sum_{x=0}^{\infty} \pi(x) = 1 \Rightarrow \sum_{x=1}^{\infty} \frac{\tilde{p}}{\tilde{q}} \left(\frac{\tilde{p}}{\tilde{q}}\right)^{x-1} \pi(0) + \pi(0) = 1$$

$$\Rightarrow \pi(0) \left[1 + \sum_{x=1}^{\infty} \left(\frac{\tilde{p}}{\tilde{q}}\right)^{x-1} \right] = 1$$

$$\Rightarrow \pi(0) = \frac{1}{1 + \sum_{x=1}^{\infty} \left(\frac{\tilde{p}}{\tilde{q}}\right)^{x-1}} \quad \begin{array}{l} > 0 \text{ iff } \tilde{p} < \tilde{q} \\ \Rightarrow \frac{\tilde{p} = \tilde{q}}{\text{Null}} \\ \text{recurrence.} \end{array}$$
$$= \frac{1}{1 + \frac{\tilde{p}/\tilde{q}}{1 - \tilde{p}/\tilde{q}}} = 1 - \tilde{p}/\tilde{q}$$

$$\Rightarrow \pi(x) = \pi(0) \left(\frac{\tilde{p}}{\tilde{q}}\right)^x = \left(1 - \frac{\tilde{p}}{\tilde{q}}\right) \left(\frac{\tilde{p}}{\tilde{q}}\right)^x$$

Lec. 21&22 were on random walk in random environment

will post later.

Branching process.

Lec. 23

63

• A population of individuals reproduces at each time step n .

$X_n =$ # of individuals at time n

• Each individual produces a random number of offsprings, and after producing them, the individual dies! (A horrible, horrible death)

Assumptions on reproduction process

1. Each individual produces offsprings with the same probability distribution: $\exists P_0, P_1, \dots, P_k, \sum_{i=0}^{\infty} P_i = 1$, so that the individual has k children with prob. P_k .

2. Individuals reproduce independently!

• Notation: The offspring distribution.

• Individual i produces Z_i^n offsprings at time n

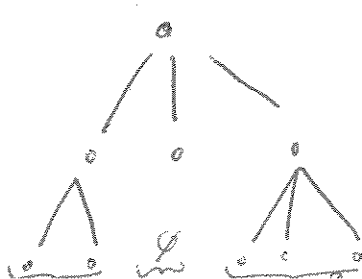
with $P\{Z_i^n = j\} = P_j \quad 0 \leq j < \infty$.

Why is it called a branching process? Example

$$X_0 = 1$$

$$X_1 = 3 = Z_1^0 + Z_2^0 + Z_3^0$$

$$X_2 = Z_1^1 + Z_2^1 + Z_3^1$$



~~attention~~ In general $X_n = \sum_{i=1}^k Z_i^n$ where

$k = X_{n-1}$. (This should make sense!)

Transition probabilities ...

Are extremely hard to compute! For example:

$$P(k, j) = \mathbb{P}\{X_{n+1} = j \mid X_n = k\} = \mathbb{P}\{Z_1^{n+1} + Z_2^{n+1} + \dots + Z_k^{n+1} = j\}$$

with Z_i^{n+1} i.i.d.

However, we can compute $p(0,0)$! ($p(0,0) = 1$! why?)

~~Ques~~ A natural set of questions!

- 1: What is the expected number of population at level n ?
- 2: ~~How can~~ Do we know if the population will go extinct?
- 3: Can we say when the population will die?

1] Let $\mu = \mathbb{E}(Z_i^n)$ = mean number of offsprings of an individual

• By definition
$$\mu = \sum_{i=0}^{\infty} i p_i$$

Then,
$$\mathbb{E}[X_{n+1} \mid X_n = k] = \mathbb{E}(Z_1^{n+1} + Z_2^{n+1} + \dots + Z_k^{n+1}) = k\mu$$

Then $E(X_n) \stackrel{!}{=} \sum_{k=0}^{\infty} P\{X_{n-1} = k\} E(X_n | X_{n-1} = k)$ (Definition)

$$= \sum_{k=0}^{\infty} k \mu P\{X_{n-1} = k\} = \mu \sum_{k=0}^{\infty} k \cdot P\{X_{n-1} = k\} \quad (*)$$

$$\stackrel{!}{=} \mu \cdot E[X_{n-1}] \quad (\text{recursion formula})$$

Iterate the procedure $\Rightarrow E[X_n] = \mu^n E[X_0]$. (*)

2 ~~Example~~

Assume $\mu < 1$. Then (*) says $E[X_n] = c \cdot \mu^n \rightarrow 0$ ($n \rightarrow \infty$).

Now observe that since by (*)

$$E[X_n] = \sum_{k=0}^{\infty} k P\{X_n = k\} \geq \sum_{k=1}^{\infty} P\{X_n = k\} = P\{X_n \geq 1\}$$

So $P\{X_n \geq 1\} \xrightarrow{n \rightarrow \infty} 0$. Thus $P\{X_n = 0\} \xrightarrow{n \rightarrow \infty} 1 - c\mu^n$

Fact: If $\mu < 1$, $\rightarrow 1$ exponentially fast.

$P\{\text{population dies eventually out}\} = 1$

Q4: What about when $\mu = 1, \mu > 1$? Intuition?

A useful tool: Probability generating function.

Lec. 24

For a r.v. X , taking values in $\{0, 1, 2, \dots\}$ the probability generating function (or generating function)

$$\phi_X(s) = \mathbb{E}(s^X) = \sum_{k=0}^{\infty} s^k \cdot \mathbb{P}\{X=k\}$$

Assume that $0 \leq s \leq 1$ so that limit always exists. ~~Assume $0 \leq s \leq 1$~~

Example:

Let X be a Poisson (λ) $\Rightarrow \mathbb{P}\{X=k\} = e^{-\lambda} \cdot \frac{\lambda^k}{k!}$

Find $\phi_X(s)$.

answer:

$$\begin{aligned} \phi_X(s) &= \mathbb{E}(s^X) = \sum_{k=0}^{\infty} s^k \cdot e^{-\lambda} \frac{\lambda^k}{k!} \\ &= e^{-\lambda} \sum_{k=0}^{\infty} \frac{(\lambda s)^k}{k!} = e^{-\lambda} e^{\lambda s} = e^{\lambda(s-1)}. \end{aligned}$$

Properties of $\phi_X(s)$

$$\mathbb{P}(X=0) \leq \phi(s) \leq \mathbb{P}(X=0) + s \sum_{k=1}^{\infty} \mathbb{P}(X=k) \leq \mathbb{P}(X=0) + s$$

$$\phi_X(0) = \mathbb{P}\{X=0\} \quad (\text{why?}) \quad \leq \phi(s) \xrightarrow{s \rightarrow 0} \mathbb{P}(X=0)$$

$$\phi_X(1) = 1 \quad (\text{why?})$$

$\phi_X(s)$ is increasing in s (why?)

Unless $\mathbb{P}(X=0)=1$,
in which case $\phi(s)=1 \forall s$

$$\hookrightarrow \phi_X'(s) = \sum_{k=1}^{\infty} k s^{k-1} \mathbb{P}\{X=k\} > 0$$

$$\phi_X''(s) = \sum_{k=2}^{\infty} k(k-1) s^{k-2} \mathbb{P}\{X=k\}$$

Now observe:

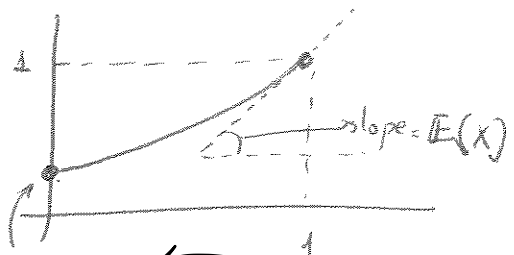
$$\phi'_X(1) = \sum_{k=1}^{\infty} k P\{X=k\} = E[X]$$

Note:
 $P(X \geq 2) = 0 \rightarrow X_n \in \{X_{n-1}, X_{n-1} + 1\}$
 Silly! $P(X=0) > 0 \Rightarrow X_n \text{ eventually } = 0$
 $P(X=1) = 1 \Rightarrow X_n = X_0 + n \Rightarrow \text{convex.}$

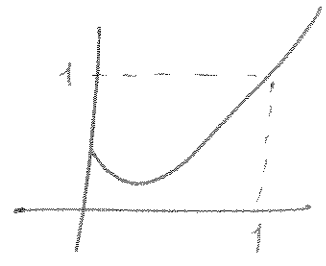
So can assume \rightarrow

$$\phi''_X(s) > 0 \iff P\{X \geq 2\} > 0.$$

General shapes



Question 5



Can it look like this?

 Not really! why?

$P\{X=0\} > 0$ w/ $X_n \rightarrow \infty$ is clear

independence: Show that $\phi_{X_1+X_2+\dots+X_m}(s) = \phi_{X_1}(s) \dots \phi_{X_m}(s)$ X_i ind

Why is $\phi_X(s)$ useful?

For branching processes: Assumptions: $p_0 > 0$; $p_0 + p_1 < 1$

(why are these assumptions reasonable?)

Define: $\alpha_n = P\{X_n = 0 \mid X_0 = 1\} = P\{\text{we are dead by time } n\}$

$$\alpha = P\{\exists n \text{ s.t. } X_n = 0 \mid X_0 = 1\}$$

$$\alpha = \lim_{n \rightarrow \infty} \alpha_n$$

Now, if the population has k individuals at a given time the only way to die out in the future is if all branches die out.

Then, ~~α^k~~ because branches behave independently
 Prob. of extinction is α^k

For all analysis purposes it suffices to determine α

Definition $\alpha = \alpha(1)$ is called the extinction probability.

Silly computation (But extremely important)

$$\begin{aligned}
a &= \mathbb{P}\{\text{population eventually dies out} \mid X_0 = 1\} \\
&= \sum_{k=0}^{\infty} \mathbb{P}\{X_1 = k \mid X_0 = 1\} \cdot \mathbb{P}\{\text{population dies out eventually} \mid X_1 = k\} \\
&= \sum_{k=0}^{\infty} p_k \alpha^k = \phi_Z(\alpha)
\end{aligned}$$

Thus, the extinction prob is a fixed point of $\phi_Z(s)$. i.e.

$$\alpha = \phi_Z(\alpha)$$

Note: $1 = \phi_Z(1)$ for all Z r.v. ^{offspring distributions}, but there might be other solutions

Assume $X_0 = 1$. Let $\phi_Z(s)$ be the offspring distribution prob. g. f.

Define $g_n(s)$ to be the generating function of X_n .

$$g_0(s) = E(s^{X_0}) = s \cdot \mathbb{P}\{X_0 = 1\} = s$$

$$g_1(s) = E(s^{X_1}) = \sum_{k=0}^{\infty} s^k \cdot \mathbb{P}\{X_1 = k\} = \sum_{k=0}^{\infty} s^k \cdot \mathbb{P}\{Z = k\} = \phi_Z(s)$$

Proposition.

$$g_n(s) = \underbrace{\phi_z \circ \phi_z \circ \dots \circ \phi_z}_{n \text{ -times}}(s) = \phi_z^{(n)}(s)$$

composition!

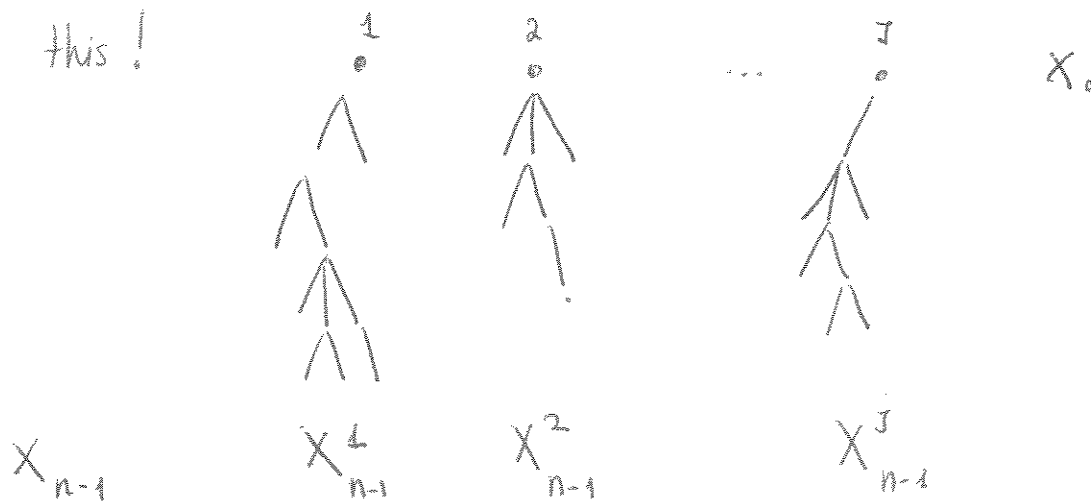
(69)

Proof.

By induction, with $n=1,0$ as base cases.

$$\begin{aligned} g_n(s) &= \sum_{k=0}^{\infty} P\{X_n=k\} s^k \\ &= \sum_{k=0}^{\infty} \sum_{j=0}^{\infty} P\{X_1=j\} \cdot P\{X_n=k | X_1=j\} s^k \\ &\stackrel{!}{=} \sum_{j=0}^{\infty} P_j \sum_{k=0}^{\infty} s^k P\{X_{n-1}=k | X_0=j\} \\ &= \sum_{j=0}^{\infty} P_j \left(\sum_{k=0}^{\infty} s^k P\{X_{n-1}=k | X_0=j\} \right) \quad (+) \\ &\quad \underbrace{\left[g_{n-1}(s) \right]^j} \end{aligned}$$

→ Proof of this!



I.I.D, (why?)

$$\Rightarrow P\{X_{n-1}=k | X_0=j\} = P\left\{ \sum_{i=1}^j X_{n-1}^i = k \mid X_0^i = 1 \right\}$$

Therefore $\sum_{k=0}^{\infty} s^k P\left\{\sum_{i=1}^j X_{n-1}^{(i)} = k\right\}$

= $g_{X_{n-1}^{(1)} + X_{n-2}^{(1)} + \dots + X_{n-2}^{(j)}}(s)$

= $g_{X_{n-1}^{(1)}}(s) g_{X_{n-1}^{(2)}}(s) \dots g_{X_{n-1}^{(j)}}(s)$

= $(g_{n-1}(s))^j$

Then (†) = $\sum_{j=0}^{\infty} p_j (g_{n-1}(s))^j \stackrel{!}{=} \sum_{j=0}^{\infty} p_j [\phi_Z^{(n-1)}(s)]^j = \phi_Z(\phi_Z^{(n-1)}(s))$
= $\phi_Z^{(n)}(s)$.
why?

Corollary (This answers $P\{X_n=0 | X_0=1\}$)

$a_n = P\{X_n=0 | X_0=1\} \stackrel{!}{=} g_n(0) = \phi^{(n)}(0)$
b/c $\phi_X(0) = P(X=0)$

Theorem

Lec. 25

The extinction probability α is the smallest possible solution

f $\alpha = \phi(\alpha)$

Proof

Let $\hat{\alpha}$ be the smallest solution such that $\hat{\alpha} = \phi(\hat{\alpha})$

Let $a_n = P\{X_n=0\}$. We'll show $a_n \leq \hat{\alpha}$

Then $a_0 = 0 < \hat{\alpha}$
 b/c $\phi(0) = p_0 \neq 0$ ϕ increasing (!)

$a_1 = P\{X_1=0\} = \phi(0) \leq \phi(\hat{\alpha}) = \hat{\alpha}$

Induction again: Assume true for $n-1$

$$a_n = P\{X_n=0\} = \phi^n(0) = \phi(\phi^{n-1}(0)) \leq \phi(\hat{a}) = \hat{a}$$

So $\alpha = \lim_{n \rightarrow \infty} a_n = P\{\text{population eventually dies out}\} \leq \hat{a}$

and $\alpha = \phi(\alpha)$ (we already know). But \hat{a} is the smallest number such that $\hat{a} = \phi(\hat{a})$ by its definition!

~~$\hat{\alpha} = \phi(\hat{\alpha})$~~ So we also have $\alpha \geq \hat{a}$. Thus: $\alpha = \hat{a}$

Question: When is $\alpha < 1$ (so we can survive?)

Theorem

if $\mu = E[Z] \leq 1$ and $p_0 > 0$, the extinction probability $\alpha = 1$

if $\mu > 1$ then $\alpha < 1$ and is the unique root (not equal to 1)

with $\phi(\alpha) = \alpha$.

Proof

If $\mu < 1$ done.

If $\mu = 1$, $\phi'_Z(1) = 1$ and $\phi''_Z > 0$. Since $\phi''_Z > 0 \Rightarrow$

$$\Rightarrow \phi'_Z \nearrow \Rightarrow \phi'_Z(s) < 1 \quad \forall s \in (0,1)$$

Then

$$1 - \phi(s) = \int_s^1 \phi'(s) ds < \int_s^1 1 ds = 1 - s \Rightarrow \boxed{s < \phi(s)} \quad \forall s \in (0,1)$$

if $\mu > 1 \Rightarrow \phi'_Z(1) > 1 \Rightarrow \exists s_0: \phi(s_0) < s_0$ (why?)

so near $s=1$ $\phi \nearrow$ but $\phi(1)=0$ so a little before $s=1$, ϕ must be < 0 .

But $\phi(0) > 0$ so the function $\phi(s)$ crosses the diagonal for $s < 1$.

some $\alpha \in (0, s_0)$

| | | | | | |
|-------|----------------------|----------------------|---------------------|--------------------------|--------------|
| D_0 | $p_0 = \frac{1}{4},$ | $p_1 = \frac{1}{4},$ | $p_2 = \frac{1}{2}$ | $\mu = 1\frac{1}{4} > 1$ | $\alpha < 1$ |
| | $p_0 = \frac{1}{2},$ | $p_1 = \frac{1}{4}$ | $p_2 = \frac{1}{4}$ | $\mu = \frac{3}{4} < 1$ | $\alpha = 1$ |
| | $p_0 = \frac{1}{4},$ | $p_1 = \frac{1}{2}$ | $p_2 = \frac{1}{4}$ | $\mu = 1$ | $\alpha = 1$ |

(72)

In the first example: α solves $\phi(\alpha) = \alpha$

$$\phi(\alpha) = \alpha^0 p_0 + \alpha^1 p_1 + \alpha^2 p_2 = \frac{1}{4} + \frac{\alpha}{4} + \frac{\alpha^2}{2}$$

$$\text{So: } \alpha = \frac{1}{4} + \frac{\alpha}{4} + \frac{\alpha^2}{2} \Rightarrow 2\alpha^2 - 3\alpha + 1 = 0$$

$$\Rightarrow (\alpha - 1)(2\alpha - 1) = 0$$

$$\text{So } \alpha = 1 \text{ or } \alpha = \frac{1}{2}$$

α is the smallest solution.

So $\alpha = \frac{1}{2}$ 50% chance of survival.